

Maryam Al-Towailb (KSU)

Discrete Mathematics and Its Applications

Math. 151 - Math. 1101

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Discrete Mathematics and Its Applications

Introduction to Logic

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Outline

- Propositional Logic
- Interpretent the second sec
- Intersection 10 The truth table for the conjunction
- The truth table for the disjunction
- The truth table for the exclusive
- Intersection of the conditional statement
- O The truth table for the biconditional statement
- S Converse, Contrapositive and inverse
- Truth tables of compound propositions
- Propositional equivalences
- Logical equivalences

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- Propositional Logic
- 2 The truth table for the negation
- The truth table for the conjunction
- The truth table for the disjunction
- S The truth table for the exclusive
- O The truth table for the conditional statement
- O The truth table for the biconditional statement
- Solution Converse, Contrapositive and inverse
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- Propositional equivalences
- Logical equivalences

Definition

A proposition is a declarative sentence (that is, a sentence that declares a fact) that is either true or false, but not both.

Example

- Washington, D.C., is the capital of the United States of America
- 3+2=6
- What time is it?
- Read this carefully.
- **(a)** x + 1 = 2

Propositions 1 is true, whereas 2 is false. Sentences 3 and 4 are not propositions because they are not declarative sentences. Sentence 5 is not propositions because they are neither true nor false.

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The truth value of a proposition is true, denoted by T, if it is a true proposition and false, denoted by F, if it is a false proposition.

Definition

Let *p* be a proposition. The negation of *p*, denoted by $\neg p$ (also denoted by \overline{p}), is the statement "It is not the case that *p*." The proposition $\neg p$ is read "not *p*." The truth value of the negation of *p*, is the opposite of the truth value of *p*.

Example

Find the negation of the proposition "Today is Friday" and express this in simple English. The negation is "It is not the case that today is Friday". This negation can be more simply expressed by "Today is not Friday," or "It is not Friday today."

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The truth table for the negation of a proposition p.



Let *p* and *q* be propositions. The conjunction of *p* and *q*, denoted by $p \land q$, is the proposition "*p* and *q*." The conjunction $p \land q$ is true when both *p* and *q* are true and is false otherwise.

Example

Find the conjunction of the propositions p and q where p is the proposition "Today is Friday" and q is the proposition "It is raining today." The conjunction of these propositions, $p \land q$, is the proposition "Today is Friday and it is raining today."

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Example

What is the disjunction of the propositions p and q where p is the proposition "Today is Friday" and q is the proposition "It is raining today." The disjunction of p and $q, p \lor q$, is the proposition "Today is Friday or it is raining today."

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The truth table for the negation of a proposition p.

TABLE the Conj Proposit	2 The Trunction of ions.	uth Table for Two	TABLE the Disju Proposit	3 The Tr inction of 1 ions.	uth Table f Iwo
P	9	$p \wedge q$	P	9	$p \lor q$
Т	Т	Т	Т	Т	Т
Т	F	F	Т	F	Т
F	Т	F	F	Т	Т
F	F	F	F	F	F

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Definition

Let p and q be propositions. The exclusive or of p and q, denoted by $p \oplus q$, is the proposition that is true when exactly one of p and q is true and is false otherwise.

TABLE 4The Truth Table forthe Exclusive Or of TwoPropositions.					
P	9	$p \oplus q$			
Т	Т	F			
Т	F	Т			
F	Т	Т			
F	F	F			

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P	9	$p \oplus q$			
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Т	F	Т			
F	Т	Т			
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P	9	$p \oplus q$
Т	Т	F
Т	F	Т
F	Т	Т
F	F	F

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Let p and q be propositions. The conditional statement $p \rightarrow q$ is the proposition "if p, then q." The conditional statement $p \rightarrow q$ is false when p is true and q is false, and true otherwise. In the conditional statement $p \rightarrow q$, p is called the hypothesis (or antecedent or premise) and q is called the conclusion (or consequence)

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TABLE 5 The Truth Table fothe Conditional Statement $p \rightarrow q$.					
P	9	$p \rightarrow q$			
Т	Т	Т			
Т	F	F			
F	Т	Т			
F	F	Т			

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Truth Tables o f Compound Propositions.

Converse, Contrapositive and inverse

- The proposition $q \rightarrow p$ is called the converse of $p \rightarrow q$.
- ② The contrapositive of $p \rightarrow q$ is the proposition $\neg q \rightarrow \neg p$.
- (a) The proposition $\neg p \rightarrow \neg q$ is called the inverse of $p \rightarrow q$.

We will see that of these three conditional statements formed from $p \rightarrow q$, only the contrapositive always has the same truth value as $p \rightarrow q$. When two compound propositions always have the same truth value we call them equivalent,

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Definition

Let p and q be propositions. The biconditional statement $p \leftrightarrow q$ is the proposition "p if and only if q." The biconditional statement $p \leftrightarrow q$ is true when p and q have the same truth values, and is false otherwise. Biconditional statements are also called bi-implications.

Diconditional $p \leftrightarrow q$.					
P	9	$p \leftrightarrow q$			
Т	Т	Т			
Т	F	F			
F	Т	F			
F	F	Т			

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Truth Tables o f Compound Propositions:

Example

Construct the truth table of the following compound propositions:

$$(p \vee \neg q) \to (p \wedge q)$$

P	9	79	$p \lor \neg q$	$p \wedge q$	$(p \vee \neg q) \to (p \land q)$
Т	Т	F	Т	Т	т
Т	F	Т	Т	F	F
F	Т	F	F	F	т
F	F	Т	Т	F	F

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TAB	LE7 T	he Truth	Table of (p	$\vee \neg q) \rightarrow (q)$	$p \wedge q$).
Р	9	¬q	$p \lor \neg q$	$p \wedge q$	$(p \vee \neg q) \to (p \wedge q)$
Т	Т	F	Т	Т	Т
Т	F	Т	Т	F	F
F	Т	F	F	F	Т
F	F	Т	Т	F	F

Propositional Equivalences:

Definition

A compound proposition that is always true, no matter what the truth values of the propositions that occur in it, is called a tautology. A compound proposition that is always false is called a contradiction. A compound proposition that is neither a tautology nor a contradiction is called a contingency.

Example

We can construct examples of tautologies and contradictions using just one propositional variable.

Consider the truth tables of $p \lor \neg p$ and $p \land \neg p$. Because $p \lor \neg p$ is always true, it is a tautology. Because $p \land \neg p$ is always false, it is a contradiction

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Logical Equivalences:

Definition

The compound propositions p and q are called logically equivalent if $p \leftrightarrow q$ is a tautology. The notation $p \equiv q$ denotes that p and q are logically equivalent.

In particular, the compound propositions p and q are equivalent if and only if the columns giving their truth values agree.

TABLE 2 De Morgan's
Laws.
$$\neg(p \land q) \equiv \neg p \lor \neg q$$

 $\neg(p \lor q) \equiv \neg p \land \neg q$

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Logical Equivalences:

Definition

The compound propositions p and q are called logically equivalent if $p \leftrightarrow q$ is a tautology. The notation $p \equiv q$ denotes that p and q are logically equivalent.

In particular, the compound propositions p and q are equivalent if and only if the columns giving their truth values agree.

pg
TABLE 2 De Morgan's
Laws.

$$\neg(p \land q) \equiv \neg p \lor \neg q$$

 $\neg(p \lor q) \equiv \neg p \land \neg q$

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Example

Show that $\neg(p \lor q)$ and $\neg p \land \neg q$ are logically equivalent.

TAB	LE 3 T	ruth Tables	for $\neg (p \lor q)$:	and ¬p ^	¬q.	
p	9	$p \lor q$	$\neg (p \lor q)$	-¬p	<i>¬q</i>	$\neg p \land \neg q$
Т	Т	Т	F	F	F	F
Т	F	Т	F	F	Т	F
F	Т	Т	F	Т	F	F
F	F	F	Т	Т	Т	Т

Example

Show that $\neg(p \lor q)$ and $\neg p \land \neg q$ are logically equivalent.

TAB	LE 3 1	ruth Tables	for $\neg (p \lor q)$:	and ¬p ^	<i>¬q</i> .	
p	9	$p \lor q$	$\neg (p \lor q)$	-¬p	<i>¬q</i>	$\neg p \land \neg q$
Т	Т	Т	F	F	F	F
Т	F	Т	F	F	Т	F
F	Т	Т	F	Т	F	F
F	F	F	Т	Т	Т	Т

Example

Show that $p \rightarrow q$ and $\neg p \lor q$ are logically equivalent.

TABLE 4 Truth Tables for $\neg p \lor q$ and $p \rightarrow q$.					
p	9	¬ <i>p</i>	$\neg p \lor q$	$p \rightarrow q$	
Т	Т	F	Т	Т	
Т	F	F	F	F	
F	Т	Т	Т	Т	
F	F	Т	Т	Т	

Example

Show that $p \rightarrow q$ and $\neg p \lor q$ are logically equivalent.

$\begin{bmatrix} TAB \\ p \rightarrow \end{bmatrix}$	TABLE 4 Truth Tables for $\neg p \lor q$ and $p \rightarrow q$.					
p	q	¬ <i>p</i>	$\neg p \lor q$	$p \rightarrow q$		
Т	Т	F	Т	Т		
Т	F	F	F	F		
F	Т	Т	Т	Т		
F	F	Т	Т	Т		

Equivalence	Name
$p \wedge \mathbf{T} \equiv p$ $p \vee \mathbf{F} \equiv p$	Identity laws
$p \lor \mathbf{T} \equiv \mathbf{T}$ $p \land \mathbf{F} \equiv \mathbf{F}$	Domination laws
$p \lor p \equiv p$ $p \land p \equiv p$	Idempotent laws
$\neg(\neg p) \equiv p$	Double negation law
$p \lor q \equiv q \lor p$ $p \land q \equiv q \land p$	Commutative laws
$(p \lor q) \lor r \equiv p \lor (q)$ $(p \land q) \land r \equiv p \land (q)$	(r) Associative laws (r)
$p \lor (q \land r) \equiv (p \lor q)$ $p \land (q \lor r) \equiv (p \land q)$	$ \begin{array}{c} \wedge (p \lor r) \\ \vee (p \land r) \end{array} $
$\neg (p \land q) \equiv \neg p \lor \neg q$ $\neg (p \lor q) \equiv \neg p \land \neg q$	De Morgan's laws
$p \lor (p \land q) \equiv p$ $p \land (p \lor q) \equiv p$	Absorption laws
$p \lor \neg p \equiv \mathbf{T}$ $p \land \neg p \equiv \mathbf{F}$	Negation laws

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