# Token Ring

Chapter 5

#### Token Ring

Computer network connecting devices in a closed loop with successive point-to-point links.

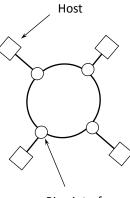
- Deterministic, asynchronous, with acknowledgment,
- Speed = 4 / 16 Mbps,
- Technology: developed by IBM,
- normalized by IEEE : 802.5,
- More complex & more expensive than the Ethernet network.

#### Token Ring

- IEEE 802.5 Standard
- Layers specified by 802.5:
  - Token Ring Physical Layer
  - Token Ring MAC Sublayer

### Token Ring (cont d)

- IEEE 802.5 Standard
- Layers specified by 802.5:
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  - Token Ring MAC Sublayer
- Token Ring, unlike Ethernet, requires an active interface.

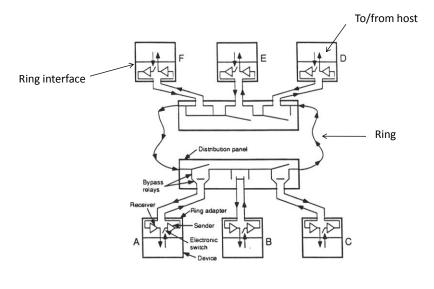


Ring interface

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### Token Ring Configuration



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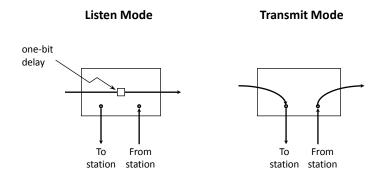
Token Ring Physical Layer

- Ring Interfaces
  - Listen and Transmit Modes
- Channel Logic
  - Differential Manchester Encoding

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### Token Ring Interface Modes





Token Ring MAC Sublayer

- Token passing protocol
- Frame format
- Token format

#### Token Passing Protocol

- A token (3 bytes pattern) circulates around the ring
- Token state:
  - Busy: supports a frame,
  - Idle: free (may be captured).

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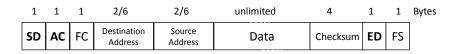
#### Token Passing Protocol (cont d)

- General Procedure:
- 1. Sending host waits for and captures an idle token,
- 2. Sending host changes the token to a frame and circulates it,
- 3. Receiving host accepts the frame and continues to circulate it,
- 4. Sending host receives its frame, removes it from the ring, and generates an idle token which it then circulates on the ring.

#### Token Ring Frame and Token Formats



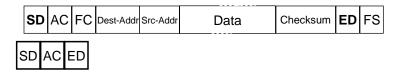
Token Format



Frame Format

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#### **Token Ring Delimiters**

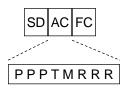


- SD = Starting Delimiter,
- ED = Ending Delimiter,
- They contains invalid differential Manchester codes.

<b>CD</b>	J	К	0	J	К	0	0	0
SD	1 bit							

J and K are code violations of the Manchester encoding and will be detected by the hardware.

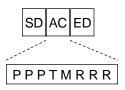
### Token Ring Access Control Field



(Note: The AC field is also used in frames)

- **P** = Priority bits (current priority)
  - provides up to 8 levels of priority when accessing the ring
- T = Token bit
  - T=0: Token
  - T=1: Frame

#### Token Ring Access Control Field (cont d)

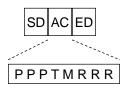


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**M** = Monitor Bit

- · Prevents tokens and frames from circulating indefinitely,
- All frames and tokens are issued with M=0,
- On passing through the "monitor station" M is set to 1,
- All other stations repeat this bit as set,
- A token or frame that reaches the monitor station with M=1 is considered invalid and is purged.

#### Token Ring Access Control Fields (cont'd)



- R = Reservation Bits
  - Allows stations with high priority data to request (in frames and tokens as they are repeated) that the next token be issued at the requested priority.



### Token Ring Frame Control Field



- FC = Frame Control Field
  - Defines the type of frame being sent,
  - Frames may be either data frames or some type of control frame. Example control frames:
    - Beacon: Used to locate breaks in the ring,
    - Duplicate address test: Used to test if two stations have the same address.

#### Token Ring Address & Data Fields

	Destination	Source	Data	Checksum	ED	FS
SUAC	Address	Address				

#### Address Fields:

- Indicate the source and destination hosts
- Broadcast:
  Set all destination address bits to 1s.
- Data
  - No fixed limit on length
  - Caveat: Hosts may only hold the token for a limited amount of time (10 msec)

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#### Token Ring Checksum and Frame Status



- Checksum: 32-bit CRC
- FS = Frame Status
  - Contains two bits, A and C
  - When the message arrives at the destination, it sets A=1
  - When the destination copies the data in the message, it sets C=1

**A** = 1 , Address recognized **C** = 1 , Frame copied

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#### The Token Ring Monitor Station

- One station on the ring is designated as the "monitor station"
- The monitor station:
  - marks the M bit in frames and tokens,
  - removes marked frames and tokens from the ring,
  - watches for missing tokens and generates new ones after a timeout period.

#### Using Priority in Token Ring

If a host wants to send data of priority *n*, it may only grab a token with priority value *n* or lower.

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- A host may reserve a token of priority *n* by marking setting the reservation bits in the AC field of a passing token or frame.
  - Caveat: The host may not make the reservation if the token or frame's AC field already indicates a higher priority reservation.
- The next token generated will have a priority equal to the reserved priority.

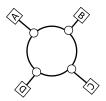
#### Using Priority in Token Ring (cont'd)

- When a new token is generated (i.e., when a sender finishes sending and releases an idle token), or when a sender sends a data frame,
  - R bits are set to the lowest priority,
  - M bit is reset.

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#### Priority Transmission: Example



Host B has 1 frame of priority 3 to send to A Host C has 1 frame of priority 2 to send to A Host D has 1 frame of priority 4 to send to A Token starts at host A with priority 0 and circulates clockwise Host C is the monitor station (priority 0: lowest priority in this example)

# Example (cont d)



Event	Token/Frame AC Field
A generates a token	P=0, M=0, T=0, R=0
B grabs the token and sets the	
message destination to A	P=3, M=0, T=1, R=0
Frame arrives at C, and C reserves	
priority level 2. Monitor bit set.	P=3, M=1, T=1, R=2
Frame arrives at D, and	
D attempts to reserve priority level 4:	P=3, M=1, T=1, R=4
Frame arrives at A, and A	
copies it	P=3, M=1, T=1, R=4
Frame returns to B, so B removes	
it, and generates a new token	P=4, M=0, T=0, R=0
Token arrives at C, but its priority is	
too high. C reserves priority 2. sets M bit.	P=4, M=1, T=0, R=2

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# Example (cont'd)

Event	Token/Frame AC Field
Token arrives at D, and D grabs	
it, sending a message to A	P=4, M=0, T=1, R=0
Frame arrives at A, and A	
copies it	P=4, M=0, T=1, R=0
Frame arrives at B, which does	
nothing to it	P=4, M=0, T=1, R=0
Frame arrives at C, which sets the	
monitor bit, C reserves priority 2	P=4, M=1, T=1, R=2
Frame returns to D, so D removes	
it and generates a new token with P=2	P=2, M=0, T=0, R=0

etc... Attempt to complete this scenario on your own.